An Overview of Aspect and Component Models: Part 3

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Outline

Component Models

- Kmelia Component Model [5 Papers]
- Fractal Component Model [1 Paper + 1 Technical Report]
- Kmelia and Fractal Models: Comparison
- Aspect Oriented Programming Issues
 - Formal Semantics for Aspects : CASB [1 Technical Report]
 - Aspect Classification [1 paper]
 - Aspect Interaction [2 paper]
 - Formalizing Concurrent Aspects [1 paper]
- Perspectives

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Kmelia is a Formal Component Model [Attiog06]

Component
$$\stackrel{def}{=} \langle W, Init, A, N, I, D_s, v, C_s \rangle$$

• $W \stackrel{def}{=} \langle T, V, V_T, Inv \rangle$
• $D_s \stackrel{def}{=} \langle I_s, B_s \rangle$
• $I_s \stackrel{def}{=} \langle \sigma, P, Q, V_s, S_s \rangle$
• $S_s \stackrel{def}{=} \langle sub_s, cal_s, req_s, int_s \rangle$
• $\mathcal{B}_s \stackrel{def}{=} \langle S, L, \delta, \phi, S_0, S_F \rangle$

Assembly $\stackrel{def}{=} < C$, links, subs >

A component composition is defined as a well-formed assembly which is encapsulated within a component.

Kmelia is a Hierarchical Component Model [Pascal06]

- Services in Kmelia are not simple operations
- Kmelia introduces the concept of Assemblies
- Kmelia proposes three hierarchy levels :
 - Links Hierarchy
 - Services interfaces Hierarchy
 - Component Composition is an encapsulation of an assembly

Kmelia defines Components Protocols [Pascal07a]

- A protocol is a pre-ordering of services calls that should be respected during the system execution.
- A protocol has a behavior
- A protocol in Kmelia is a specific service defined using vertical structuring operators
 - Sate annotation << >>
 - Transition annotation [[]]
- Protocol inconsistency detection can be made using pre/post conditions.

Kmelia introduces HBIDL to describe components and services [Pascal07b]

- HBIDL extends IDL by the specification of the behavior of services with their architectures
- HBIDL has many advantages:
 - provides detailed documentations of complex interaction services
 - supports compatibility levels
 - serves as an intermediate between CBSE and SBSE
- HBIDL has some adaptation problems such as:
 - Parameters vs Messages mismatch
 - Hierarchichal mismatch

Kmelia has a Formal Anlyser Toolbox: COSTO [Pascal07c]

- COSTO is a toolbox that supports the design and analysis of Kmelia's abstract component model
- COSTO is an eclipse plugin
- COSTO toolbox includes:
 - COSTO core module
 - Verification module
 - LOTOS Module
 - MEC Module
 - Export Module
- COSTO takle state explosion problem

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Fractal Component Model (1) [Bruneton04, Bruneton06]

A Fractal Component is an entity that has two parts:

- Membrane
- Content

Fractal Model supports three kind of Components:

- Basic Components
- Primitive Components
- Composite Components

Fractal Supports two kinds of Components Binding:

- Primitive binding
- Composite Binding

Fractal Component Model (2) [Bruneton04, Bruneton06]

Fractal Component Model has the following main features:

- Fractal is a hierarchical model
- Fractal supports sharing components
- Fractal is a reflective model
- Fractal has an implementation model named Julia

CORBA Component Model (CCM) [OMG'04, Marvie'06]



- $CCM = \sum Component Types + \sum Component Homes + \sum Links$
 - Component Type = Name + Attributes + [Ports]
 - Attributes are used for configuration purpose
 - Ports = Facets \cup Receptacles \cup Event Sources \cup Event Sinks
 - An inheritance relationship is defined between component types
 - Component Types are two kinds : Basic and Extended
 - Component Home is a meta-type that manages component instances

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CCM: Facets [OMG'04, Marvie'06]

- A Facet defines a role that can be performed upon a component
- Each Facet has its own reference
- Several copies of a Facet reference may exist at a time.
- Facets can be specified only in a static way
- A client can navigate between Facets

CCM: Receptacles [OMG'04, Marvie'06]

- A receptacle allows a component type to accept a reference
- A receptacles maybe simples or multiples
- A cookie is created for each connection in the case of multiple receptacles
- Receptacles can be used for reconfiguration.

CCM: Events [OMG'04, Marvie'06]

- Events are used for asynchronous communications
- Event sources are two kinds: Emitters and Publishers
- An Event Sink may receive events from various sources at the same time

CCM: Component Homme [OMG'04, Marvie'06]

- A Component Home is a component manager that provides instantiation of component types at runtime
- Many home types may manage the same component type only with different instances
- Component Homes are two kinds: Keyless and with primary key Homes
- Component Homes are not components \Rightarrow non hierarchical model

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CCM: Components Configuration [OMG'04, Marvie'06]

- A Component configuration is implemented using configuration objects
- A Component home has Factory operations for component instances

Components Models

Aspect Oriented Programming Issues

Perspectives

CCM: Global Software Production Process [OMG'04, Marvie'06]



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Kmelia and Fractal Component Models: A Comparison

- Kmelia is Service Based Model ≠ Fractal is a Component Based Model
- Kmelia follows monadic semantics ≠ Fractal follows demi-polyadic semantics
- Three hierarchy levels are allowed in Kmelia ≠ One hierarchy level for Fractal
- No sharing Components for Kmelia ≠ Sharing Components is allowed with Fractal
- reconfiguration is limited in Kmelia ≠ reconfiguration is more developed in Fractal

Components Models

Aspect Oriented Programming Issues

Perspectives

Formal Semantics for Aspects

CASB: Semantic Elements

- CASB introduces the concept of configurations (C , Σ)
- A program C is of the form $C ::= i : C | \varepsilon$
- Semantic is described in term of binary relation \rightarrow_b
- A single reduction : $(i : C, \Sigma) \rightarrow_b (C', \Sigma')$
- An aspect is a function
 ψ : I → (Σ → C) × {before, after, around}
- A tagged instruction : \overline{i}
- A matching function: $match : \mathcal{P} \times I \rightarrow bool$
- A weaving relation: \rightarrow

CASB: Weaving of a Single Aspect

• Before aspect

$$\frac{\psi(i) = (\phi, before)}{(i:C, \Sigma) \to (test \ \phi: \overline{i}:C, \Sigma)}$$

• After aspect

$$\frac{\psi(i) = (\phi, after)}{(i: C, \Sigma) \to (\overline{i}: test \ \phi: C, \Sigma)}$$

• Around aspect

$$\frac{\psi(i) = (\phi, around)}{(i:C, \Sigma, P) \to (test \ \phi: pop_p: C, \Sigma, \overline{i}:P)}$$
$$(pop_p:C, \Sigma, i:P) \to (C, \Sigma, P)$$
$$(proceed:C, \Sigma, i:P) \to (i: push_p \ i:C, \Sigma, P)$$

CASB: Weaving of Several Aspects

- Aspects of the same Kind
 - Before aspects

$$\frac{\psi(i) = ((\phi_1 ... \phi_n), before)}{(i: C, \Sigma) \to (test \ \phi_1 : ... : test \ \phi_n : \overline{i} : C, \Sigma)}$$

• After aspects

$$\frac{\psi(i) = ((\phi_1 \dots \phi_n), after)}{(i: C, \Sigma) \to (\overline{i}: test \ \phi_1 : \dots : test \ \phi_n : C, \Sigma)}$$

• Around aspects

$$\frac{\psi(i) = ((\phi_1 \dots \phi_n), around)}{(i: C, \Sigma, P) \to (test \ \phi_1 : pop_p \ n : C, \Sigma, test \ \phi_2 : \dots : test \ \phi_n : \overline{i} : C, \Sigma)}$$
$$(pop_p \ n : C, \Sigma, x_1 : \dots : x_n : P) \to (C, \Sigma, P)$$
$$(proceed : C, \Sigma, x : P) \to (x : push_p \ x : C, \Sigma, P)$$

• Aspects of Different Kinds $\frac{\psi(i) = ((\phi_1, t_1) \dots (\phi_n, t_n)) \quad \gamma((\phi_1, t_1) \dots (\phi_n, t_n)) = ((\phi'_1, around) \dots (\phi'_n, around))}{(i:C, \Sigma, P) \rightarrow (test \ \phi'_1: pop_p \ n:C, \Sigma, test \ \phi'_2: \dots: test \ \phi'_n: \vec{i}:P)}$

CASB: Pointcuts

$$\mathcal{P} ::= T_i | \mathcal{P}_1 \land \mathcal{P}_2 | \mathcal{P}_1 \lor \mathcal{P}_2 | \neg \mathcal{P}$$

$$match(T_i, i) = true \text{ if } \exists \sigma : \sigma(T_i) = i$$

$$= \text{ false otherwise}$$

$$match(P_1 \land P_2, i) = match(P_1, i) \land match(P_2, i)$$

$$match(\neg P, i) = \neg match(P, i)$$

CASB: Exception Handling

• Exception Syntax :

$$S ::= try S_1 catch ex S_2 | throw ex | \dots$$

• Semantics :

$$(try S_1 \ catch \ ex \ S_2 : C, \Sigma, E) \rightarrow_b (S_1 : pop_e : C, \Sigma, (ex, S_2 : C) : E)$$

(throw $ex : C, \Sigma, (ex_0, C_0) : \ldots : (ex_k, C_k) : (ex, C') : E) \rightarrow_b (C', \Sigma, E)$

$$(pop_e: C, \Sigma, X: E) \rightarrow_b (C, \Sigma, E)$$

CASB: Advanced Aspect Features

• Aspect Deployment

$$(deploy id S: C, \Sigma, \Psi) \to (S: pop_{\Psi}: C, \Sigma, \psi_{id}: \Psi)$$

$$(pop_{\Psi}: C, \Sigma, \psi_i: \Psi) \to (C, \Sigma, \Psi)$$

• Aspect Instantiation

$$\frac{update(\Psi, i, \Sigma) = (\Psi', \Sigma') \ (\circ\Psi')(i) = (\phi, before)}{(i, C, \Sigma, \Psi) \to (test \ \phi : \overline{i} : C, \Sigma', \Psi')}$$

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Aspects Classification

- Organize aspects into categories sharing some properties
- -- > Specify the preserved properties by the aspects of each category
- --> Optimization of the verification time

Observers Category

• Definition :

$$\forall (C, \Sigma) . \Sigma^{\psi} \in \mathcal{A}_a \Leftrightarrow proj_b(\alpha) = proj_b(\tilde{\alpha}) \land preserve_b(\tilde{\alpha})$$

• Preserved Properties :

$$\begin{array}{lll} \varphi^{o} & ::= & sp \mid \neg sp \mid \varphi_{1}^{o} \land \varphi_{2}^{o} \mid \varphi_{1}^{o} \lor \varphi_{2}^{o} \\ & \mid \varphi_{1}^{o} \cup \varphi_{2}^{o} \mid \varphi_{1}^{o} W \varphi_{2}^{o} \mid true \cup \varphi^{'o} \\ \varphi^{'o} & ::= & ep \mid \neg ep \mid sp \mid \neg sp \mid \varphi_{1}^{'o} \land \varphi_{2}^{'o} \mid \varphi_{1}^{'o} \lor \varphi_{2}^{'o} \\ & \mid \varphi_{1}^{'o} \cup \varphi_{2}^{'o} \mid \varphi_{1}^{'o} W \varphi_{2}^{'o} \mid true \cup \varphi^{'o} \end{array}$$

Aborters Category

- Definition : $\forall (C, \Sigma). \Sigma^{\psi} \in \mathcal{A}_o \Leftrightarrow (proj_b(\alpha) = proj_b(\tilde{\alpha}) \lor \exists (i \geq 0), \exists (j \geq i). proj_b(\alpha_{\rightarrow i}) = proj_b(\tilde{\alpha}_{\rightarrow j}) \land \forall (k > j). \tilde{\alpha}_k = (\epsilon, ..)) \land preserve_b(\tilde{\alpha})$
- Preserved Properties :

$$\begin{array}{lll} \varphi^{a} & ::= & sp \mid \neg sp \mid \varphi^{a}_{1} \land \varphi^{a}_{2} \mid \varphi^{a}_{1} \lor \varphi^{a}_{2} \mid \varphi^{a}_{1} W \varphi^{a}_{2} \mid true \cup \varphi^{'a} \\ \varphi^{'a} & ::= & \neg ep \mid \varphi^{'a}_{1} \land \varphi^{'a}_{2} \mid \varphi^{'a}_{1} \lor \varphi^{'a}_{2} \mid true \cup \varphi^{'o} \end{array}$$

Aspect Interactions: Detection and Resolution (EAOP)[Rémi'02]

• An aspect in EAOP is :

$$A ::= \mu a.A$$
$$| C \triangleright I; A$$
$$| C \triangleright I; a$$
$$| A_1 \Box A_2$$

• Crosscuts and Inserts :

$$C ::= T \mid C_1 \wedge C_2 \mid C_1 \vee C_2 \mid \neg C$$

• A term T is :

$$T ::= f T_1 \dots T_n \mid x$$

• Aspect Composition :

$$(\mu a. C_1 \triangleright I_1; a) || (\mu a. C_2 \triangleright I_2; a)$$

Aspect Weaving (1)

$$sel j (\mu a.A) = sel j A$$

$$sel j (C \triangleright I; A) = \phi \qquad \text{if C j=fail}$$

$$= \{C. \triangleright I\} \qquad \text{otherwise}$$

$$sel j (A_1 \Box A_2) = sel j A_1 \qquad \text{if sel } j A_1 \neq \phi$$

$$= sel j A_2 \qquad \text{otherwise}$$

Aspect Weaving (2)

• The Monitor:

$$[j, P, \sigma]^{\phi} \mapsto \sigma[end]$$

$$\frac{\mathcal{S} = \{C \triangleright I\} \cup \mathcal{S}' \quad C \ j = \psi \quad (\downarrow, \psi I, \sigma) \xrightarrow{*} (\uparrow, \psi I, \sigma')}{[j, P, \sigma]^{\mathcal{S}} \mapsto [j, P, \sigma']^{\mathcal{S}'}}$$

• Woven Execution:

$$\frac{[j, P, \sigma]^{\operatorname{sel} jA} \stackrel{*}{\mapsto} \sigma_a \quad (j, P, \sigma_a) \to (j', P, \sigma')}{(A, j, P, \sigma) \Rightarrow (\operatorname{next} jA, j', P, \sigma')}$$

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Aspect Strong Independence

• Laws for Aspects :

[(un)fold] $\mu a.A = A[\mu a.A/a]$ $(A_1 \Box A_2) \Box A_3 = A_1 \Box (A_2 \Box A_3)$ [assoc] $(C_1 \triangleright I_1; A_1) \Box (C_2 \triangleright I_2; A_2) = (C_2 \triangleright I_2; A_2) \Box (C_1 \triangleright I_1; A_1)$ [commut] if $C_1 \wedge C_2 = fail$ if C = failelim₁ $C \triangleright I = false \triangleright I$ elim₂ $(false \triangleright I; A_1) \Box A_2 = A_2$ elim false $\triangleright I$; $C_1 \triangleright I_1$; $A = false \triangleright I$; A $(C_1 \triangleright I_1; A_1) \Box (C_2 \triangleright I_2; A_2) = (C_1 \triangleright I_1; A_1) \Box (C_2 \land \neg C_1 \triangleright I_2; A_2)$ [priority] let $A = (C_1 \triangleright I_1; A_1) \Box \dots \Box (C_n \triangleright I_n; A_n)$ [propag] and $A' = (C'_1 \triangleright I'_1; A'_1) \Box \dots \Box (C'_m \triangleright I'_m; A'_m)$ then $A \parallel A' = \Box_{i=1\dots n}^{j=1\dots m} C_i \wedge C'_i \triangleright (I_i \bowtie I'_i); (A_i \parallel A'_i)$ $\Box_{i=1..n}C_i \triangleright I_i; (A_i \parallel A')$ $\Box_{j=1..m}C'_j \triangleright I'_j; (A \parallel A'_j)$

Aspect Independence w.r.t a Program

$$I_{w}(A,j) = if \text{ sel } jA = \phi \text{ then } \Box_{j' \in (\text{step}_{p} j)} I_{w}(A,j')$$

$$else if \text{ sel } jA = \{C \triangleright I\} \text{ then } C \triangleright I; \ \Box_{j' \in (\text{step}_{p} j)} I_{w}(\text{next } jA,j')$$

$$else if \text{ sel } jA = \{C \triangleright I, \ C' \triangleright I'\} \text{ then } C \wedge C' \triangleright (I \bowtie I'); \ \Box_{j' \in (\text{step}_{p} j)} I_{w}(\text{next } jA,j')$$

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Conflict Resolution

- Using parallel operators such us $\|_{seq}$ and $\|_{fst}$
- Defining scopes for Aspects: scope id Idset A

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Extended version for the Framework [Rémi'04]

• Inter-crosscut Variables :

$$C ::= v \doteq T | C_1 \land C_2 | C_1 \lor C_2 | \neg C$$

$$C = C[(\bullet \doteq T/\hat{T})]$$

$$C = C[z/\bar{x}] \land x \doteq z$$
next $j (C \triangleright I; A) = C \triangleright I; A \text{ if } C \text{ j=fail}$

$$= \psi A \text{ otherwise}$$

$$\psi(\mu a.A) = \mu a.\psi A$$

$$\psi(C \triangleright I; A) = \psi' C \triangleright \psi' I; \psi' A$$

$$\psi(A_1 \Box A_2) = (\psi A_1 \Box \psi A_2)$$

Composition Operators

- Sequential Composition Operator : $A_1 C \rightarrow A_2$
- Adaptors $A_1 \parallel_O A_2$:

$$O ::= \mu a.O | C \triangleright F; O | C \triangleright F; a | O_1 \Box O_2 F ::= (U \oplus B) U ::= id | skip B ::= \bowtie | seq | fst | snd | skip$$

Aspect Requirements

- Requirements are aspects
- CFA is expressed by:

$$J ::= f J_1...J_n | ?$$

$$S ::= \mu s.S \\ | (J_1 \to S_1)[]...[](J_n \to S_n) \\ | s$$

Sequential EAOP [Rémi'06]

- Aspects are translated from EAOP syntax to the FSP
- Events are introduced to synchronize advices together or advices with the base program
- Weaving is modeled as parallel composition of the FSP describing the base program and the Aspect.

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Concurrent EAOP

- Each program is viewed as a parallel composition of several FSPs automatons describing its threads
- Aspects are viewed as independent processes that run in parallel and synchronized with the base program
- A hiding operator is introduced to implement the concurrency.

Concurrent Aspect Composition in EAOP

- Sequential Functional Composition: *Fun*(*aspect*₁, *aspect*₂)
- Parallel Conjunctive Composition: : *parAnd*(*aspect*₁, *aspect*₂)

Components	Models
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Perspectives