

An Overview of CoCoME

Hakim Hannousse



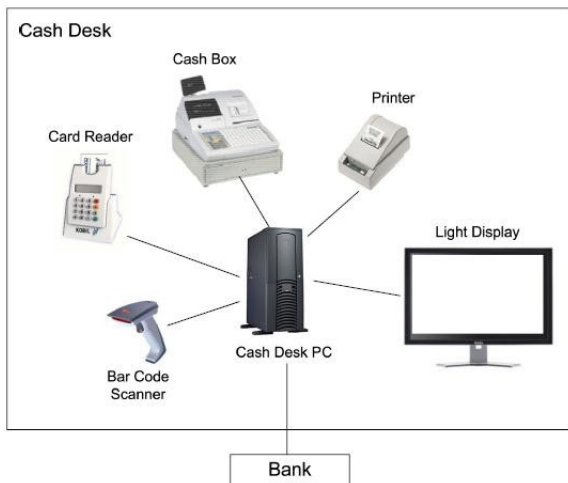
ECOLE DES MINES DE NANTES

DEPARTMENT OF COMPUTER SCIENCE

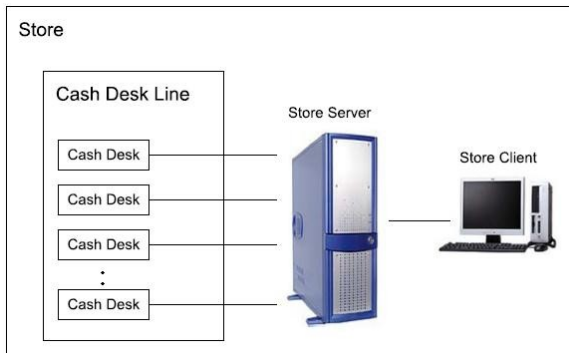
CoCoME: Common Component Modeling Example

- Context = Trading System
- Supports functional aspects : Manage sales, order products, .. etc.
- Supports non function aspects : Manage Express Checkout, Synchronization, RealTime Constraints ... etc.
- Extra-functional properties based on statistics for typical German super-markets

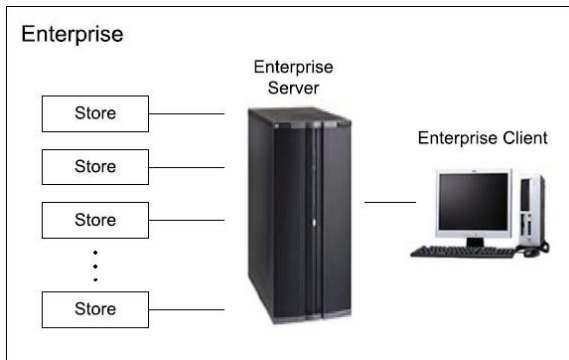
CoCoME: A Single Cash Desk



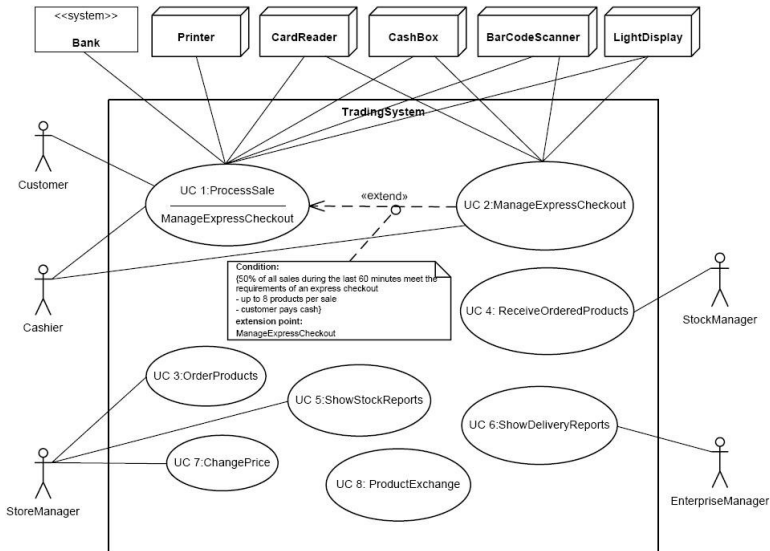
CoCoME: A single Store



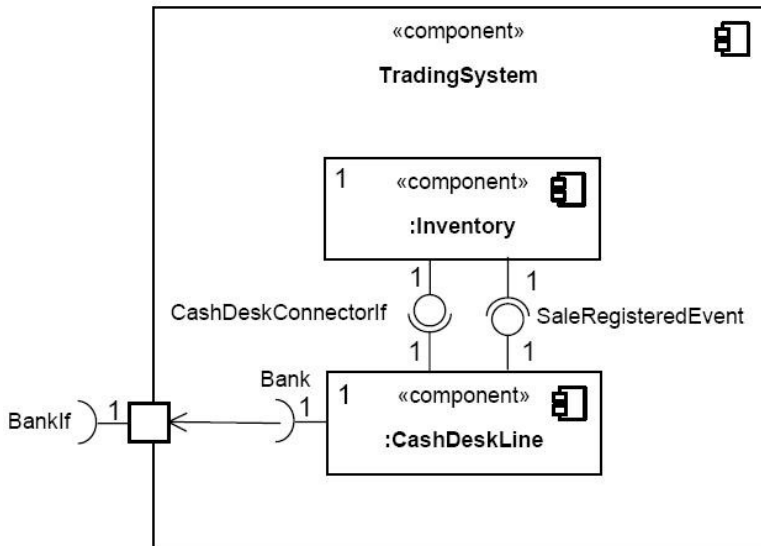
CoCoME: An Enterprise



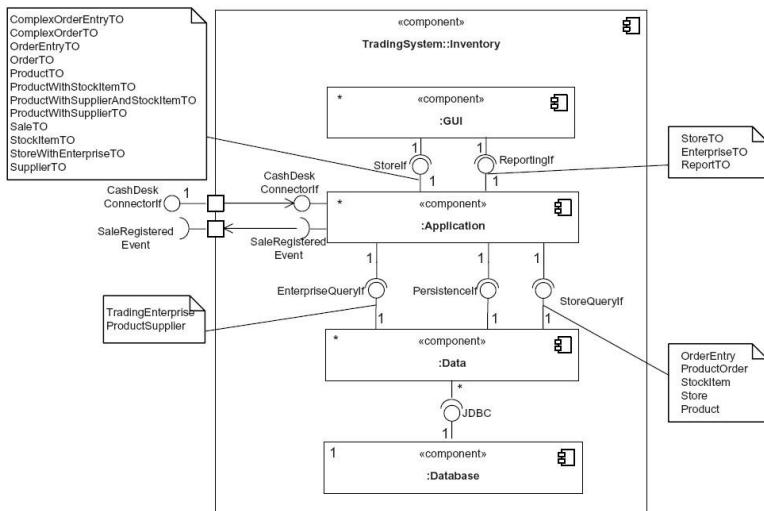
CoCoME: Provided Use Cases



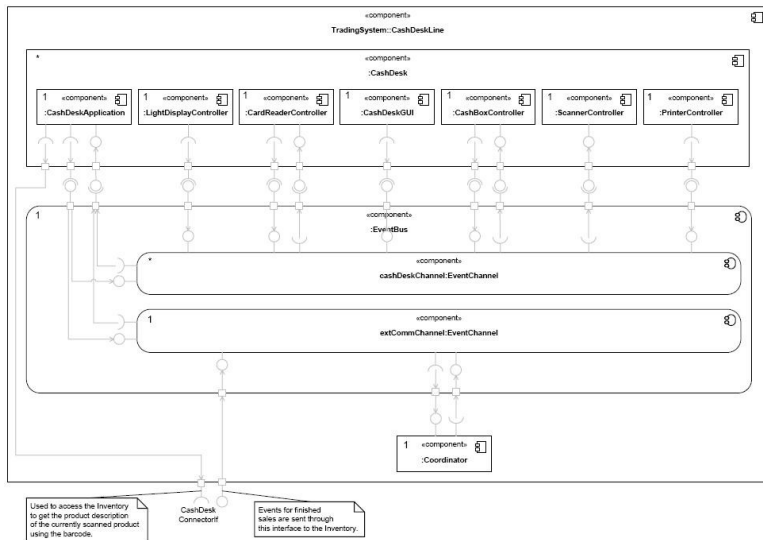
CoCoME: Component Modularization - TradingSystem



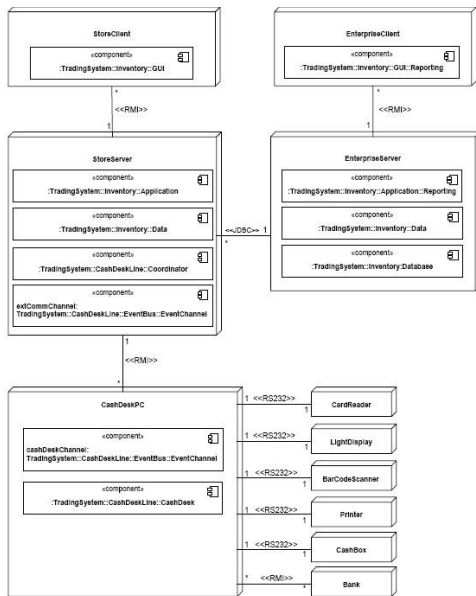
CoCoME: Component Modularization - Inventory



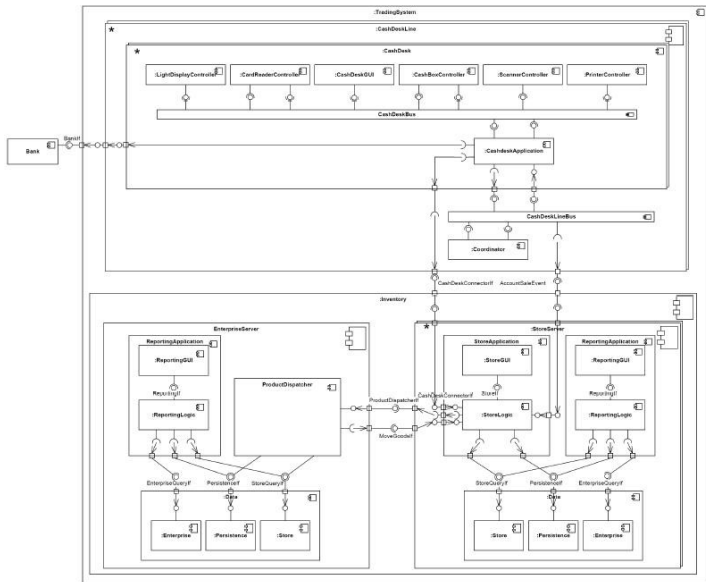
CoCoME: Component Modularization - CashDeskLine



CoCoME: Component Modularization - Deployment System



CoCoME in Fractal : Structural View



CoCoME in Fractal : Behavioral View

- ① A frame protocol is associated to each component
- ② Frame protocol language used by FractalBPC :

$$\begin{array}{l}
 P ::= !I.M\uparrow \\
 \quad | ?I.M\downarrow \\
 \quad | !I.M\downarrow \\
 \quad | ?I.M\uparrow \\
 \quad | ?I.M\{P\} \\
 \quad | !I.M\{P\} \\
 \quad | P+ \\
 \quad | P* \\
 \quad | P_1|P_2 \\
 \quad | P_1;P_2
 \end{array}$$

CoCoME in Fractal : Behavioral View

```
# INITIALISED
(
  ?CashDeskApplicationHandler.onSaleStarted
);

# SALE_STARTED
(
  ?CashDeskApplicationHandler.onProductBarcodeScanned{
    !CashDeskConnector.getProductWithStockItem;
    !CashDeskApplicationDispatcher.sendProductBarcodeNotValid+
    !CashDeskApplicationDispatcher.sendRunningTotalChanged
  }
)*; # <--- LOOP

?CashDeskApplicationHandler.onSaleFinished;

# SALE_FINISHED
(
  ?CashDeskApplicationHandler.onPaymentMode
);

# PAYING_BY_CASH
(
  (
    (
      ?CashDeskApplicationHandler.onCashAmountEntered
    )*)
  )
);
```

CoCoME in Fractal : Behavioral View

```
# On Enter
?CashDeskApplicationHandler.onCashAmountCompleted{
    !CashDeskApplicationDispatcher.sendChangeAmountCalculated
};

?CashDeskApplicationHandler.onCashBoxClosed{
    !CashDeskApplicationDispatcher.sendSaleSuccess;
    !CDLEventDispatcher.sendAccountSale;
    !CDLEventDispatcher.sendSaleRegistered
}
)
)
)* | (
# Enable Express Mode
?CDLEventHandler.onExpressModeEnabled{
    !CashDeskApplicationDispatcher.sendExpressModeEnabled
}
)* | (
# Disable Express Mode
?CashDeskApplicationHandler.onExpressModeDisabled
)*
```

CoCoME in Fractal : Deployment View

- ① FractalRMI are used rather than Sun RMI
- ② JMS are not used for implementing buses, they are replaced by components routing messages
- ③ Deployment is described using FractalADL and implemented using FractalRMI

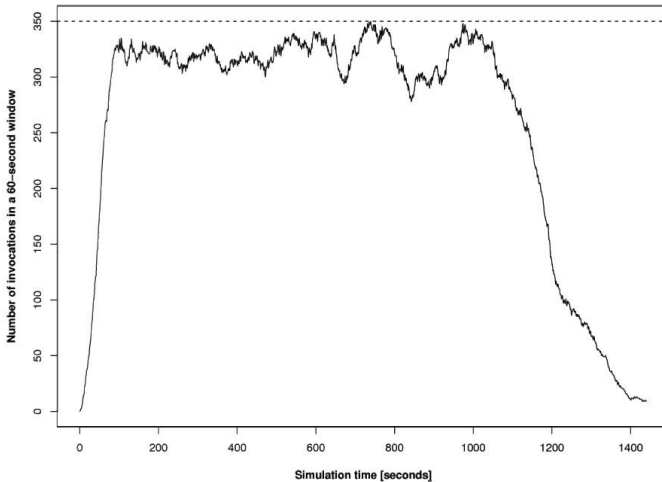
CoCoME in Fractal : Implementation View

- 1 The architecture is modeled using FractalGUI
- 2 The resulting model is extended by hand to integrate behavior protocols
- 3 A tool is used to get the skeleton of the application
- 4 The code of the CoCoME implementation is adapted and inserted to the corresponding components

CoCoME in Fractal : Testing process

- FractalBPC is used to check components communicates behaviors
 - ① The GUI components are not considered for testing
 - ② Extra-functional proprieties are independently tested from the functionality testing
 - ③ The trading system is automatically lunched.

CoCoME in Fractal : Testing Results



SOFA is a Hierarchical Component Model

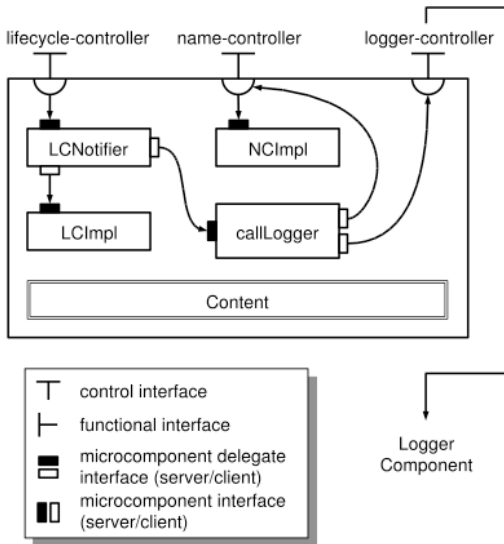
- Each Component is defined by:
 - Frame : provided and required interfaces
 - Architecture : subcomponents and their interconnections
- A Component has two parts :
 - Control part
 - Content part
- Components are bound using connectors

SOFA is ADL-Based design

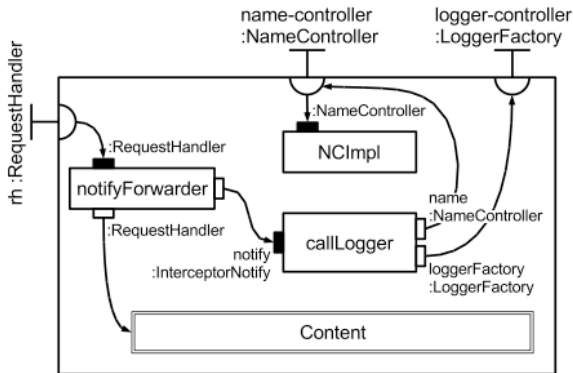
```
<frame name="TradingSystemFrame">  
  <requires name="BankIf" itf-type="BankIf"/>  
</frame>
```

```
<architecture name="TradingSystemArch" frame="TradingSystemFrame">  
  <sub-comp name="CashDeskLine" frame="CashDeskLineFrame"  
arch="CashDeskLineArch"/>  
  <sub-comp name="Inventory" frame="InventoryFrame" arch="InventoryArch"  
>  
  <connection>  
    <endpoint itf="AccountSaleEventHandlerIf" sub-comp="CashDeskLine"/>  
    <endpoint itf="AccountSaleEventHandlerIf" sub-comp="Inventory"/>  
  </connection>  
  <connection>  
    <endpoint itf="CashDeskConnectorIf" sub-comp="CashDeskLine"/>  
    <endpoint itf="CashDeskConnectorIf" sub-comp="Inventory"/>  
  </connection>  
  <connection>  
    <endpoint itf="BankIf" sub-comp="CashDeskLine"/>  
    <endpoint itf="BankIf"/>  
  </connection>  
</architecture>
```

SOFA 2.0 is A Microcomponent-Based Component Controller Model (1)



SOFA 2.0 is A Microcomponent-Based Component Controller Model (2)



SOFA 2.0 Supports Controllers Extensions Using Aspects

```

<aspect-definition name="logging" >
  <frame-addon-definition name="logger-itfs" >
    <interface signature="LoggerFactory"
      role="client" name="logger-controller" />
  </frame-addon-definition>
  <component name="logger"
    definition="logger-adl" />

  <microcomponent-definition
    name="callLogger" >
    <interface signature="InterceptorNotify"
      role="server" name="notify" />
    <interface signature="LoggerFactory"
      role="client" name="loggerFactory" />
    <interface signature="NameController"
      role="client" name="name" />
    <content class="LoggerInterceptor" />
  </microcomponent-definition>

  <microcomponent-definition
    name="notifyForwarder" >
    <interface signature="InterceptorNotify"
      role="client" name="notify" />
    <dynamic-interface role="delegateserver" />
    <dynamic-interface role="delegateclient" />
    <content
      generator="InterceptorNotifyGenerator" />
  </microcomponent-definition>

```


SOFA 2.0 Supports Controllers Extensions Using Aspects

```
<select-component type="any" >

  <frame-addon definition="logger-itsf" />
  <component-binding client="this.logger-controller"
    server="logger.logFactory" />

  <select-interface name="*" type="functional">

    <microcomponent name="logFwd"
      definition="notifyForwarder"
      flow="passthrough" />
    <microcomponent name="logCalls"
      definition="callLogger" flow="standalone"/>

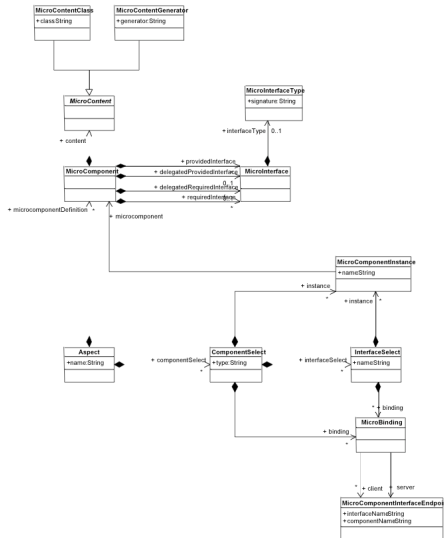
    <binding client="logCalls.loggerFactory"
      server="this.logger-controller" />
    <binding client="logCalls.name"
      server="this.name-controller" />
    <binding client="logFwd.notify"
      server="logCalls.notify" />

  </select-interface>
</select-component>
</aspect-definition>
```

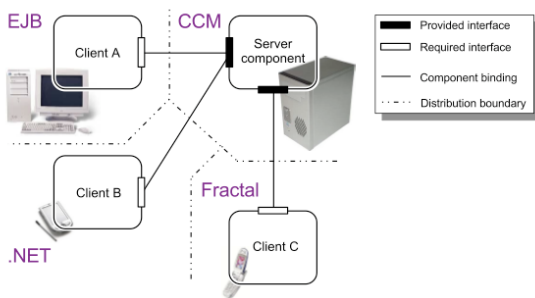
SOFA 2.0 Supports Controllers Extensions Using Aspects

```
<configuration>
  <aspect name="protocols" definition="...protocols" />
  <aspect name="logging" definition="...logging" />
  <apply-aspect name="protocols" />
</aspect>
<application definition="examples.hello.Hello">
  <apply-aspect name="protocols" />
  <apply-aspect name="logging">
    <param name="logger-key" value="auditlog" />
    <target path="./server" />
  </apply-aspect>
</application>
</configuration>
```

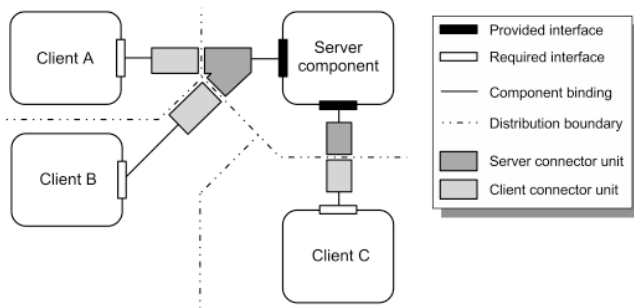
SOFA 2.0 - Microcomponent Meta-Model



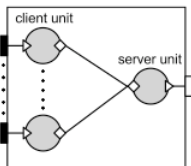
SOFA 2.0 Supports Heterogeneous Deployment via First-Entity Class Connectors (1)



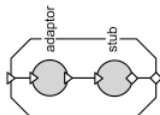
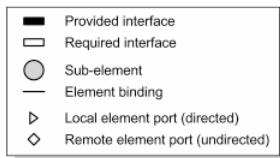
SOFA 2.0 Supports Heterogeneous Deployment via First-Entity Class Connectors (2)



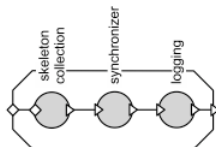
SOFA 2.0 - Connector Architecture



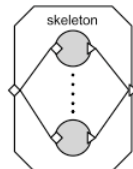
a) connector architecture



b) client unit element architecture



c) server unit element architecture



d) skeleton collection element architecture

SOFA 2.0 - Formal Signature of ports on an element

RMI Skeleton

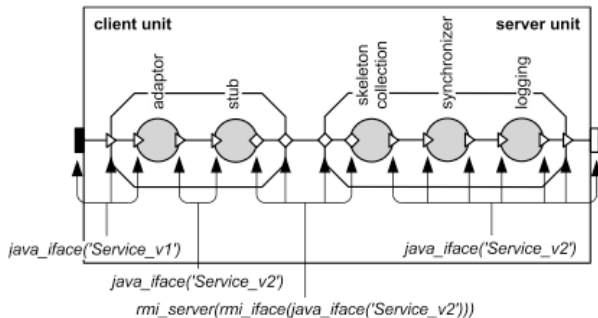


Port signatures:

call: *ltf*

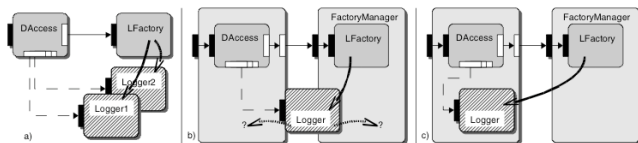
line: *rmi_server(rmi_iface(ltf))*

SOFA 2.0 - Interface adaptation and propagation through a connector



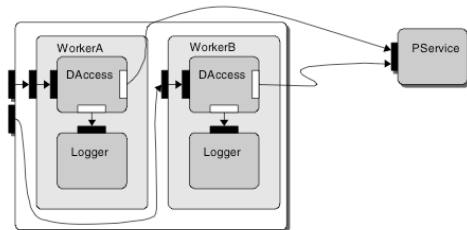
SOFA 2.0 - Supports Dynamic Reconfiguration Only w.r.t Configuration Patterns

1 Nested Factory Pattern

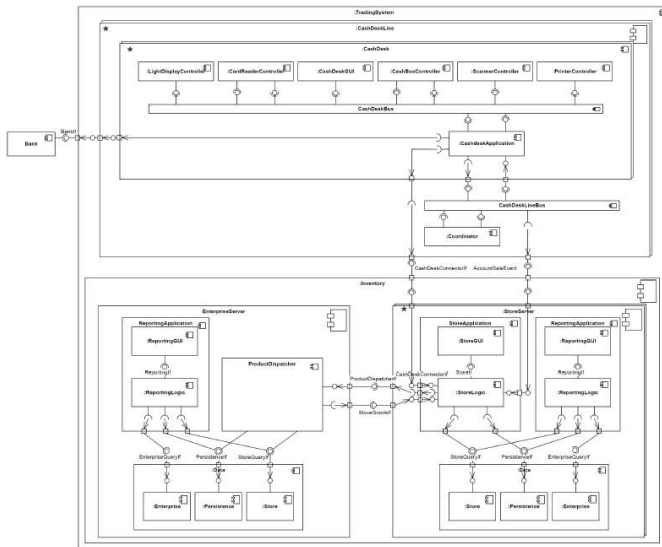


2 Removing Component Pattern

3 Utility Interface Pattern



SOFA 2.0 - Structural View



SOFA 2.0 - Behavioral View

SOFA 2.0 uses EBP to describe component's behavior :

- Example :
(*?i.open; (?i.read+?i.write)*; ?i.close*)|*?ctrl.status**
- EBP supports types and local variable declarations
- EBP supports switch and while statements

```
component LightDisplay {  
  
  types {  
    states = {LIGHT_ENABLED, LIGHT_DISABLED}  
  }  
  
  vars {  
    states state = LIGHT_ENABLED  
  }  
  
  behavior {  
    ?LDispCtrlEventHandlerIf.onEvent(EVENT ExpModeEnabledEvent){  
      state <- LIGHT_ENABLED  
    }*  
    |  
    ?LDispCtrlEventHandlerIf.onEvent(EVENT ExpModeDisableEvent){  
      state <- LIGHT_DISABLED  
    }*  
  }  
}
```

SOFA 2.0 - Deployment View

- SOFANode distributed runtime environment is used for the deployment issue.
- SOFANode = Repository + deployment docks
- Connectors encapsulate middleware and support different communication style.
- SOFA Application lifecycle:
 - ① Defining primitive components or frame components by the developer
 - ② Uploads them in the repository
 - ③ Assembly process to construct component architectures
 - ④ A deployer assigns components to docks, sets components properties values and the control aspects to be applied in the applications in the deployment plan
 - ⑤ Connectors are generated automatically
 - ⑥ Launch the application

SOFA 2.0 - Verification and Analysis

- Compliance both vertical and Horizontal via Promela. (EBP2PR)

Table 1. The result of vertical compliance verification of the CashDesk component

# of states	EBP2PR [s]	Verification [s]	Total time [s]
3 335 950	41.5	46.1	95,6

- Verification of Code against Frame Protocols via JPF tool.
- Runtime Checking against Code
- Performance Analysis

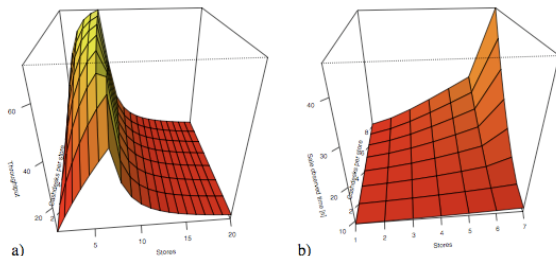


Fig. 4. Calculated a) throughput and b) average service time of Use Case 1

rCOS Component Model : Component Definition

- A Component is defined by its interface and its contract :

- Interface : $I = \langle FDecl, MDecl \rangle$

- Contract : $Ctr = \langle I, Init, MSpec, Prot \rangle$

- Example : Buffer Component

$$I_{Buffer} = \langle buff : seq(Int), \{put(inx : int), get(outy : int)\} \rangle$$

$$Ctr_{Buffer} = \langle I_{Buffer}, Init_{Buffer}, MSpec_{Buffer}, Prot_{Buffer} \rangle$$

$$Init_{Buffer} = |buff| = 0,$$

$$MSpec_{Buffer}(put(inx : int)) = (\vdash buff' = \langle x \rangle \bullet buff),$$

$$MSpec_{Buffer}(get(outy : int)) = (\vdash buff' = tail(buff) \wedge y' = head(buff)),$$

$$Prot_{Buffer} = (put, get) * +(put; (get; put)*)$$

- A Component in rCOS is a tuple:

$$C = \langle I, Init, MCode, PriMDec, PriMCode, InMDec \rangle$$

rCOS Component Model : Component Composition Operators

● Chaining operator :

Definition 2. Let C_1 and C_2 be components such that $C_1.I.FDec \cap C_2.I.FDec = \emptyset$, $C_1.I.MDec \cap C_2.I.MDec = \emptyset$ and $C_1.PriMDec \cap C_2.PriMDec = \emptyset$. Then the chaining C_1 to C_2 , denoted by $C_1 \rangle C_2$, is the component with

- $(C_1 \rangle C_2).FDec \stackrel{def}{=} C_1.FDec \cup C_2.FDec$,
- $(C_1 \rangle C_2).InMDec \stackrel{def}{=} (C_2.InMDec \cup C_1.InMDec) - (C_2.MDec \cup C_1.MDec)$,
- $(C_1 \rangle C_2).MDec \stackrel{def}{=} C_1.MDec \cup C_2.MDec$,
- $(C_1 \rangle C_2).Init \stackrel{def}{=} C_1.Init \wedge C_2.Init$,
- $(C_1 \rangle C_2).Code \stackrel{def}{=} C_1.Code \cup C_2.Code$, and
- $(C_1 \rangle C_2).PriCode \stackrel{def}{=} C_1.PriCode \cup C_2.PriCode$.

Example :

$$\begin{aligned}
 C_1.FDec &= \{buff_1:Seq(int)\} \\
 C_1.MDec &= \{put(\mathbf{in} \ x:int), get_1(\mathbf{out} \ y:int)\} \\
 C_1.Code(put) &= (buff_1 := \langle x \rangle) \triangleleft buff_1 = \langle \rangle \triangleright (put_1(\mathbf{head}(buff_1)); buff_1 := \langle x \rangle) \\
 C_1.Code(get_1) &= (buff_1 \neq \langle \rangle) \longrightarrow (y := \mathbf{head}(buff_1); buff_1 = \langle \rangle) \\
 C_1.InMDec &= \{put_1(\mathbf{in} \ x:int)\} \\
 \\
 C_2.FDec &= \{buff_2:Seq(int)\} \\
 C_2.MDec &= \{put_1(\mathbf{in} \ x:int), get(\mathbf{out} \ y:int)\} \\
 C_2.Code(put_1) &= (buff_2 = \langle \rangle) \longrightarrow buff_2 := \langle x \rangle \\
 C_2.Code(get) &= (y := \mathbf{head}(buff_2); buff_2 = \langle \rangle) \triangleleft buff_2 \neq \langle \rangle \triangleright get_1(y) \\
 C_2.InMDec &= \{get_1(\mathbf{in} \ y:int)\}
 \end{aligned}$$

rCOS Component Model : Component Composition Operators

- Disjoint Composition :

Definition 15. (Disjoint Composition) *Let C_1 and C_2 be components such that they do not share fields, public operations. Then $C_1 \otimes C_2$ is defined to be the composite component which has the provided operations of C_1 and C_2 as its provided operations, and the required operations of C_1 and C_2 as its required operations:*

$$(C_1 \otimes C_2)(InCtr) \stackrel{def}{=} C_1(InCtr|C_1.InMDec) || C_2(InCtr|C_2.InMDec)$$

rCOS Component Model : Component Composition Operators

- Feedback :

Definition 16. (Feedback) Let C be a component and $m \in C.MDec$ and $n \in C.InMDec$. $C[m \hookrightarrow n]$ is the component such that for any $InCrt$

$$C[m \hookrightarrow n](InCrt) \stackrel{def}{=} C(InCrt.MSPec \oplus \{n \mapsto (g \& [c])\}) \setminus \{m\}$$

$C.MCode(m) = g \longrightarrow c$. Notice here the design $[c]$ is the weakest fixed point of a recursive equation if it calls other methods [15].

rCOS Component Model : Component Composition

● Component Interaction Compatibility :

Definition 17. (Interaction compatibility) For a provided protocol $PProt_1$ and a required protocol $RProt_2$ given in the previous paragraph, we say they are compatible if $PProt_1 \setminus InMDec_2 \supseteq RProt_2[?op/!op \mid op \in InMDec]$, where a sequence in the required protocol is of the form $\langle !op_1(x_1), \dots, !op_k(x_k) \rangle$ and $!op_i(x_i)$ is the call out event i to operation op .

Furthermore, when they are compatible, we define the (largest) provided protocol after the provided operations are plugged in the required operations

$$PProt_1 \setminus RProt_2 \stackrel{def}{=} PProt_1 / RProt_2$$