

N-Party Rendezvous versus Sending Messages

Jean-Claude Royer¹

¹École des Mines de Nantes
Department of Computer Science – OBASCO Group
INRIA Research Centre Rennes - Bretagne Atlantique – LINA

20/09/2008

Outline

- 1 Lamport Example
- 2 2-party versus messages
- 3 N-party versus 2-party

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- Try to better understand the use of N-party rendezvous
- To compare its applicability with binary messages
- Not directly related to extraction of components from code
- However could be used to compare the models we have
- Kmelia uses a mutliparty, STSLib has N-party rendezvous, ...

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Binary Message versus Rendezvous

- **Binary message is the main way for components to communicate**
 - It is a synchronisation which triggers an action on the receiver side
 - Rendezvous principle as in Ada is the same
 - We rather want to analyse general rendezvous as in LOTOS or BIP
 - 2-party rendezvous can synchronously execute an action on each side, one for the emitter and one for the receiver
 - It can be viewed as an atomic and related set of binary messages

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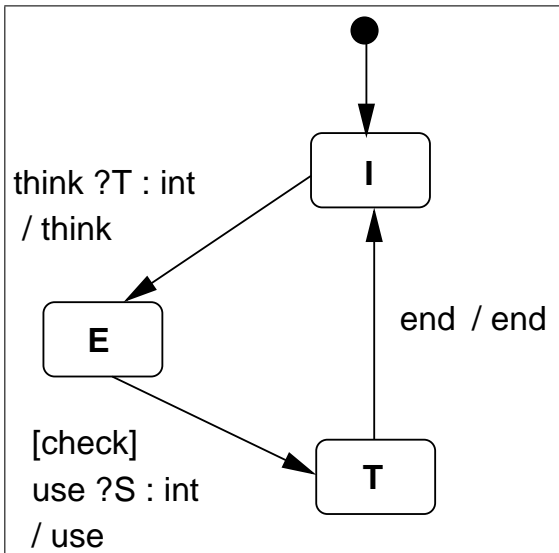
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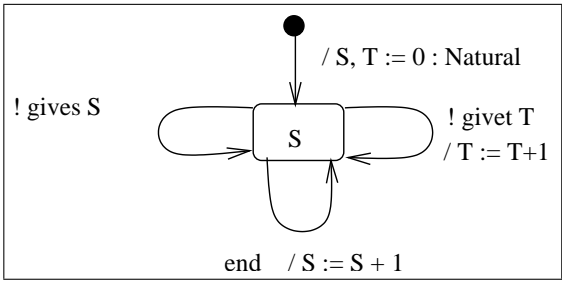
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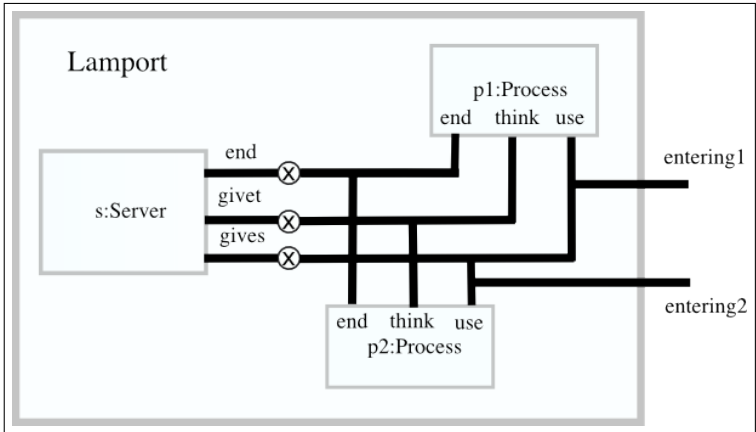


The server component



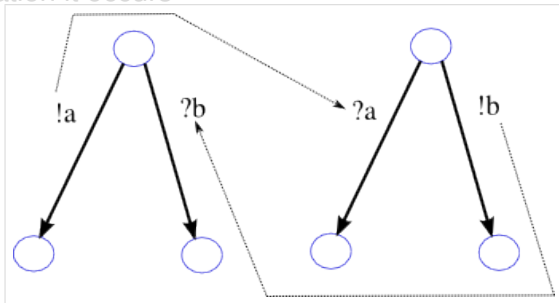
An Architecture with Two processes

Lamport
Example
2-party versus
messages
N-party
versus 2-party



Mixed State problem

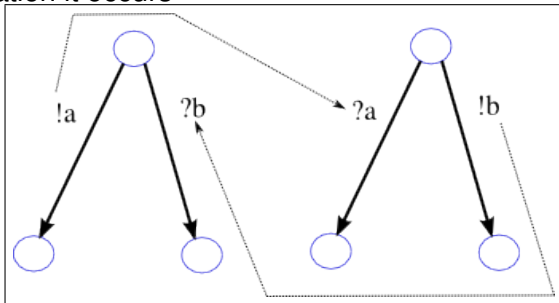
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- With simple messages and without a global coordination it occurs



- An STS message is a message without mixed state problem

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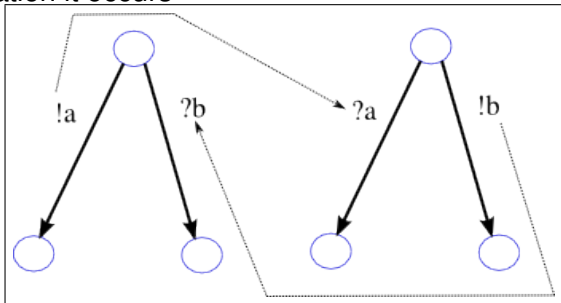
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Some Transformations

- 2-party rendezvous can be transformed into set of messages
- Double action: `server.givet / T++ - process.think / A=T`
- Guard with receipt: `server.gives !S - process.[A=S] use ? S`
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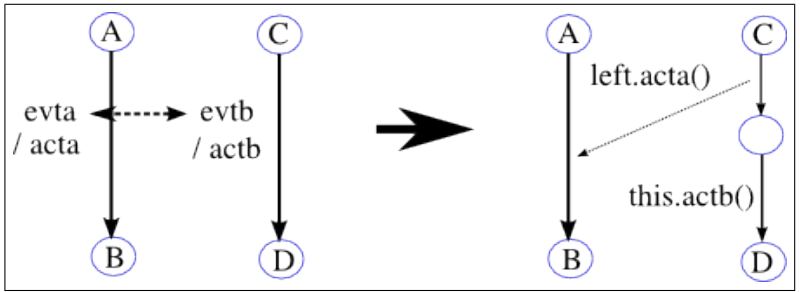
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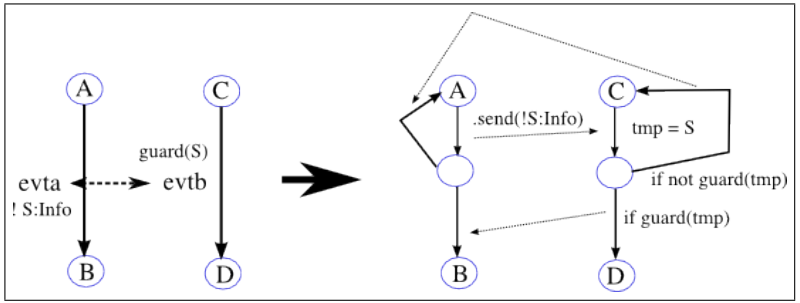
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Double Action Transformation



Guard with Receipt Transformation

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Main principles but

- **Additional constraints come from ! and ?**
- Guards with receipt imply more events, more actions and more synchronisations
- There are also some possible variations with
 - message direction
 - ordering of actions
- Transformations increase the complexity (number of states and transitions)
- Double adds 1+1 and Guard with receipt 2+4

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- Design as a 2-party system, it works with any number of processes
- Change it into a binary message system
- Change one double action the `givet` - `think` interaction
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- With four processes: `product *3 2` and `cfg *2.7 1.8`
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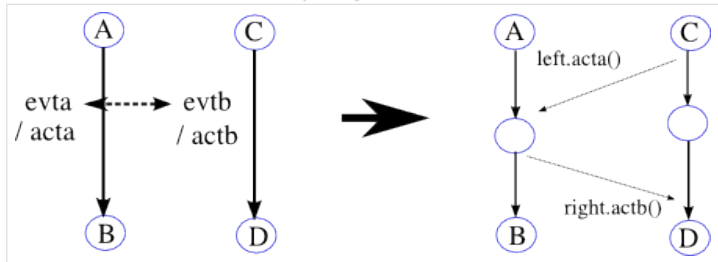
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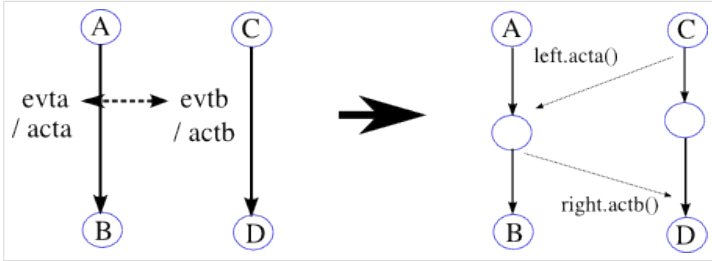
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- Cannot be generalised to N without adding more glue code

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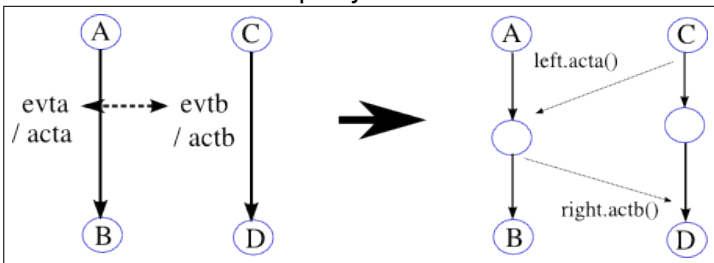
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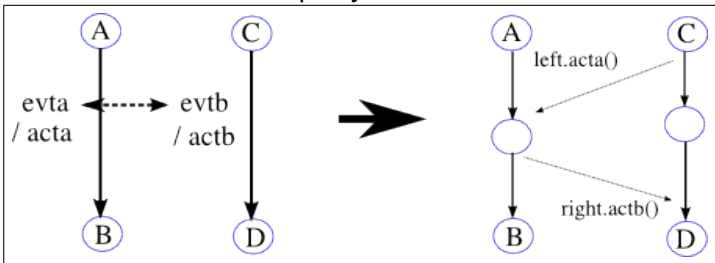
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Message with return

- Not existing with N-party, such a generalisation does not seem natural and even useful
- Even with message it does not seem a flexible way to interact with components
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N-party rendezvous

- **N-party rendezvous is a generalisation to N participants**
- It allows one way or multiway communication
- It needs two synchronisation barriers, one for entering the synchronisation area and one for leaving the area.
- With distributed systems it is not realistic
- However in local network it is possible to use it
- Automation systems, controller synthesis, modular robotics, ...
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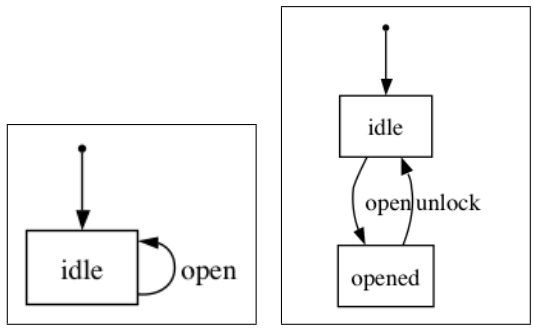
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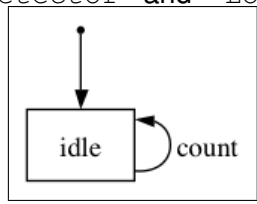
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The Primitive Components

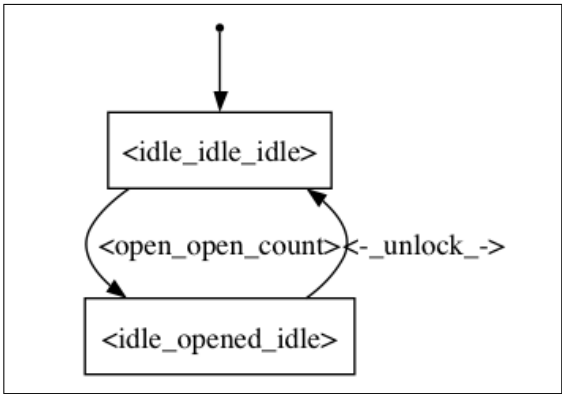
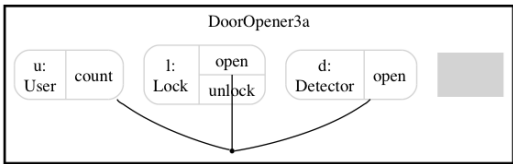


Detector and Lock



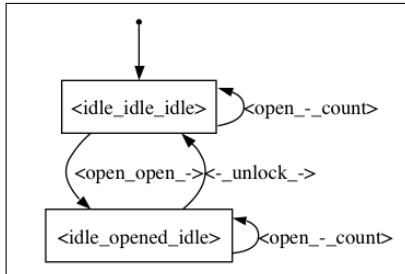
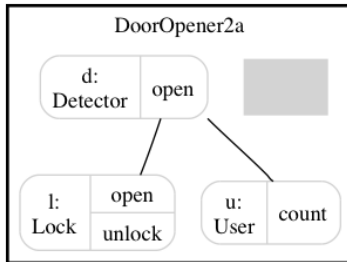
User

DoorOpener with 3-party



Lampport
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DoorOpener with 2-party

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- Not a proof but may be not too difficult to see
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Summary

- There are strictly more opportunity to compound with N-party: there are more combinations and some of them cannot be realised with binary interactions
- Powerful, but costly, mechanism
- Local network
- More abstract thus simpler behaviour
- On the other hand: The use of messages
 - Increase complexity and add problems
 - Add adaptors or controllers thus more or less a N-party mechanism
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